

## Ramsey-like theorems

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ABSTRACT: Ramsey-like theorems. Hales-Jewett.

In this note, a **coloring** is an edge-coloring unless specified otherwise.

**Entrance.** (jk: Jeudi, 26 mars, 1985) Let  $K_n$  denote the *complete graph* on  $n$  vertices. Let  $\mathcal{R}(n)$  denote the  $n^{\text{th}}$  **Ramsey number**, that is, the smallest  $M$  such that if one colors each edge of  $K_M$  either **blue** or **green**, then there must exist a monochromatic subgraph isomorphic to  $K_n$ .

**Notation.** For vertex  $\mathbf{w}$  in a **blue-green** graph, let  $bV(\mathbf{w})$  be the *set* of vertices connected to  $\mathbf{w}$  by a **blue** edge; use  $gV(\mathbf{w})$  for those connected to  $\mathbf{w}$  by **green** edges. We thus have vertex degrees

$$\text{bDeg}(\mathbf{w}) := |bV(\mathbf{w})| \quad \text{and} \quad \text{gDeg}(\mathbf{w}) := |gV(\mathbf{w})|.$$

$$\text{Hence } \text{bDeg}(\mathbf{w}) + \text{gDeg}(\mathbf{w}) = \text{Deg}(\mathbf{w}).$$

**1a: Theorem.** For  $n = 1, 2, 3, \dots$ ,

$$1a': \quad \mathcal{R}(n+1) \leq 2[1 + [n-1]\mathcal{R}(n)]. \quad \diamond$$

**Proof.** With  $M := \text{RhS}(1a')$ , let  $\mathbb{V}$  denote the vertex set of  $K_M$ . FTSOC, assume we have edge-colored  $K_M$  so that it has *no* monochromatic copy of  $K_{n+1}$ .

Pick a  $\mathbf{u}_0 \in \mathbb{V}$ . WLOG generality at least half of the edges from  $\mathbf{u}_0$  are **blue**. Letting  $B_0 := \text{bDeg}(\mathbf{u}_0)$ , we have

$$*: \quad B_0 \geq \left\lceil \frac{M-1}{2} \right\rceil \stackrel{\text{note}}{=} 1 + [n-1] \cdot \mathcal{R}(n).$$

**Building a blue complete-subgraph.** Pick a vertex  $\mathbf{u}_1 \in bV(\mathbf{u}_0)$ . The set  $bV(\mathbf{u}_0) \setminus \{\mathbf{u}_1\}$  can have at most  $\mathcal{R}(n) - 1$  vertices not belonging to  $bV(\mathbf{u}_1)$ . If otherwise, then there would be a copy of  $K_{\mathcal{R}(n)}$  in our graph, disjoint from  $\mathbf{u}_0$  and  $\mathbf{u}_1$ , such that each vertex in this copy was connected to  $\mathbf{u}_0$  by a **blue** edge and to  $\mathbf{u}_1$  by a **green** edge. This copy would, by hypothesis, contain a monochromatic copy –call it  $H$ – of  $K_n$ . Were  $H$  **blue**, then  $H \sqcup \{\mathbf{u}_0\}$  is a **blue**  $K_{n+1}$ . Else,  $H$  is **green**, so  $H \sqcup \{\mathbf{u}_1\}$  is a **green**  $K_{n+1}$ . Either is a  $\bowtie$ .

Consequently the set  $bV(\mathbf{u}_0) \cap bV(\mathbf{u}_1)$  has at least

$$|bV(\mathbf{u}_0)| - |\{\mathbf{u}_1\}| - [\mathcal{R}(n) - 1] \stackrel{\text{note}}{=} B_0 - \mathcal{R}(n)$$

many vertices. Pick some  $\mathbf{u}_2$  in  $bV(\mathbf{u}_0) \cap bV(\mathbf{u}_1)$ . The same argument shows  $bV(\mathbf{u}_0) \cap bV(\mathbf{u}_1) \cap bV(\mathbf{u}_2)$  has at least

$$[B_0 - \mathcal{R}(n)] - \mathcal{R}(n)$$

many vertices.

Continue. At stage  $n-1$  we will have chosen distinct vertices  $\mathbf{u}_0, \mathbf{u}_1, \dots, \mathbf{u}_{n-1}$ . Further, the intersection of their **blue** vertex-sets will satisfy

$$|bV(\mathbf{u}_0) \cap bV(\mathbf{u}_1) \cap \dots \cap bV(\mathbf{u}_{n-1})| \geq B_0 - [n-1]\mathcal{R}(n)$$

and hence be non-empty, by (\*). Picking a vertex  $\mathbf{u}_n$  in this intersection, now every pair of vertices in  $\{\mathbf{u}_0, \mathbf{u}_1, \dots, \mathbf{u}_n\}$  is connected by a **blue** edge.  $\bowtie \spadesuit$

**Remark.** Values  $\mathcal{R}(1), \mathcal{R}(2)$  and  $\mathcal{R}(3)$  are 1, 2 and 6 respectively. Using the above theorem, easily

$$\mathcal{R}(n+1) \leq [n! \cdot 2^n]. \quad \square$$

**More colors.** If we allow  $\mu \in \mathbb{Z}_+$  many colors, then we get a corresponding sequence of Ramsey numbers  $\mathcal{R}_\mu(n)$  for  $n = 1, 2, \dots$ . Here,  $\mathcal{R}_1(n) = n$ , and  $\mathcal{R}_2(\cdot)$  is another name for  $\mathcal{R}(\cdot)$ .

For each  $\mu > 1$ , lump the  $\mu$  colors into two **supercolors** consisting each of at most  $h := \lceil \mu/2 \rceil$  colors. Applying  $\mathcal{R}(\cdot)$  to these two supercolors, implies that

$$\mathcal{R}_\mu(n) \leq \mathcal{R}(\mathcal{R}_h(n)) \leq \mathcal{R}(\mathcal{R}(\dots \mathcal{R}(n) \dots)),$$

where the RhS has  $\lceil \log_2(\mu) \rceil$  occurrences of  $\mathcal{R}(\cdot)$ .  $\square$

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**Infinite Ramsey's Theorem**

(jk, Feb1989) Let  $K_\infty$  denote the complete graph on denumerably many vertices.

**1b: Infinite Ramsey's Theorem.** *Each edge  $\overline{uv}$  of  $H = (\mathbb{V}, \mathbb{E})$ , a  $K_\infty$  graph, is colored either blue or green. Then  $H$  includes a monochromatic  $K_\infty$ .*  $\diamond$

**Proof.** Let  $W_0 := \mathbb{V}$ . At stage  $N$  we have vertex-sets

$$W_0 \supset W_1 \supset \dots \supset W_j \supset \dots \supset W_N,$$

each infinite. Moreover, we have vertices  $(\mathbf{u}_j)_{j=0}^{N-1}$ , with  $\mathbf{u}_j \in W_j \setminus W_{j+1}$ , so that each edge-set

$$\mathbb{E}_j := \{\overline{v u_j} \mid v \in W_{j+1}\}$$

is monochromatic.

Continue the induction by picking an arbitrary vertex  $\mathbf{u}_N \in W_N$ . Since  $W_N \setminus \{\mathbf{u}_N\}$  is infinite, it includes an infinite set of vertices  $\mathbf{v}$  so that  $\{\overline{v u_N}\}_v$  is monochromatic. Define  $W_{N+1}$  to be this infinite set of vertices.

**An infinite monochromatic sequence.** Each of the sets  $\mathbb{E}_0, \mathbb{E}_1, \dots$  is either blue or green. Thus there is a subsequence  $N_1 < N_2 < \dots$  so that, WLOG, each  $\mathbb{E}_{N_j}$  is green. The consequence is that the complete graph with vertex-set  $(\mathbf{u}_{N_j})_{j=1}^\infty$  has every edge green.  $\diamond$

**1c: Observation.** *Infinite Ramsey's Thm implies the Finite Ramsey's Thm, by a compactness argument.*  $\diamond$

**Proof.** Consider  $K_\infty$ , the complete graph with vertex-set  $\mathbb{V} = \{1, 2, \dots\}$ .

FTSOC suppose the, say, 5<sup>th</sup> Ramsey number were not finite. Then for each  $n$ , there would exist a blue-green coloring of the edges on vertex-set  $\{1, \dots, n\}$  –call this colored graph  $H_n$ – so that  $H_n$  does not include a monochromatic copy of  $K_5$ .

Interpret  $H_n$  as an edge-coloring of  $K_\infty$ : Color **plum** each edge that is not between vertices  $[1..n]$ . (I.e, for posints  $\mathbf{u} < \mathbf{v}$ : If  $\mathbf{v} > n$  then color edge  $\overline{uv}$  **plum**.)

We now have a sequence  $H_1, H_2, H_3, \dots$  of  $\{\text{blue, green, plum}\}$  colorings of  $K_\infty$ . Since each edge

of  $K_\infty$  has only finitely many [three] possible colors, the sequence of colorings has a convergent subsequence  $(H_{n_j})_{j=1}^\infty$ . The coloring obtained in the limit, an edge-coloring of  $K_\infty$ , has no **plum**. But  $\infty$ -Ramsey's Thm guarantees a monochromatic copy of  $K_\infty$ ; WLOG it is **green**. Letting  $\mathbf{u}_1 < \mathbf{u}_2 < \dots < \mathbf{u}_5$  be its first 5 vertices, this **green**  $K_5$  was already a subgraph of  $H_{\mathbf{u}_5}$ .  $\diamond$

**2a: Two-variable  $\mathcal{R}()$ .** For posints  $\mathbf{b}, \mathbf{g}$ , let  $\mathcal{R}(\mathbf{b}, \mathbf{g})$  be the smallest  $M$  st. each blue-green-coloring of  $K_M$  has either a **blue**  $K_\mathbf{b}$  or a **green**  $K_\mathbf{g}$ . So  $\mathcal{R}(n) = \mathcal{R}(n, n)$ .

Evidently,  $\mathcal{R}(\mathbf{b}, 1) = 1$  and  $\mathcal{R}(\mathbf{b}, 2) = \mathbf{b}$ ; and this holds symmetrically, since  $\mathcal{R}(\mathbf{b}, \mathbf{g}) = \mathcal{R}(\mathbf{g}, \mathbf{b})$ .  $\square$

**2b: Theorem.** *For all  $\mathbf{b}, \mathbf{g} \in \mathbb{Z}_+$ ,*

$$* \quad \mathcal{R}(\mathbf{b}, \mathbf{g}) \leq \mathcal{R}(\mathbf{b}-1, \mathbf{g}) + \mathcal{R}(\mathbf{b}, \mathbf{g}-1). \quad \diamond$$

**Proof.** Let  $M := \text{RhS}(*)$ . Fix a vertex  $\mathbf{w}$ ; it has  $M-1$  edges, so has either at least  $\mathcal{R}(\mathbf{b}-1, \mathbf{g})$  **blue**-edges, or at least  $\mathcal{R}(\mathbf{b}, \mathbf{g}-1)$  **green**-edges. WLOG

$$|\text{bV}(\mathbf{w})| \geq \mathcal{R}(\mathbf{b}-1, \mathbf{g}). \quad [\text{Recall } \text{bV}(\mathbf{w}) \text{ is the set of vertices blue-connected to } \mathbf{w}.]$$

If the graph induced by  $\text{bV}(\mathbf{w})$  has a **green**  $K_\mathbf{g}$ , then DONE. Else, the induced graph admits a **blue**  $K_{\mathbf{b}-1}$  which, together with vertex  $\mathbf{w}$ , induces a **blue**  $K_\mathbf{b}$ .  $\diamond$

**2c: Exercise:** Coloring with  $\mu$  colors, let  $\mathcal{R}(n_1, \dots, n_\mu)$  be the smallest  $M$  st. each  $\mu$ -coloring,  $H$ , of  $K_M$  has an index  $j$  for which  $H$  admits a  $K_{n_j}$  with all edges the  $j^{\text{th}}$ -color. **Exer-E1:** What is the  $\mu$ -color analog of Thm 2b? [Hint: Don't jump to conclusions.]  $\square$

**Examples from Bona's text**

In class we proved  $\mathcal{R}(3)=6$ . In this instance (2b\*) is sharp, as

$$\mathcal{R}(3,3) \leq 2 \cdot \mathcal{R}(3,2) = 2 \cdot 3 = 6.$$

Using Thm 2b again,

$$\mathcal{R}(3,4) \leq \mathcal{R}(2,4) + \mathcal{R}(3,3) = 4 + 6 = 10.$$

3a: **Claim:**  $\mathcal{R}(3,4) = 9$ . ◆

**Pf of  $\mathcal{R}(3,4) \leq 9$ .** FTSOC, suppose  $H$  is a blue-green-coloring of  $K_9$  with no blue triangle, nor green  $K_4$ .

Could *some* vertex  $w$  have  $b\text{Deg}(w) \geq 4$ ? Since no blue triangle, no pair of vertices in  $bV(w)$  has a blue-edge; but then  $bV(w)$  induces a green  $K_4$ .  $\bowtie$

Could *every*  $w$  have  $b\text{Deg}(w) = 3$ ? Then the blue-degree-sum is  $3 \cdot 9$ , which is odd. But this degree-sum must also equal twice the number of blue edges.  $\bowtie$

So *there exists* a vertex  $w$  with  $b\text{Deg}(w) \leq 2$ , hence

$$|gV(w)| \geq 8 - 2 = 6.$$

Since  $\mathcal{R}(3) = 6$ , and  $gV(w)$  cannot have a blue triangle, it must have a green triangle,  $G$ . But then  $G \sqcup \{w\}$  is a green  $K_4$ . ◆

**Pf of  $\mathcal{R}(3,4) > 8$ .** On vertex-set  $\mathbb{V} := [0..8]$ , let  $\equiv$  and  $\oplus$  and  $\ominus$  each operate mod 8.

Connect vertices  $j,k \in \mathbb{V}$  by blue IFF  $k \ominus j$  is either

$3$ , producing a blue octagon, or

$4$ , producing four center-crossing blue edges. ◆

But no triple of numbers from  $\{3, 4\}$  has  $\oplus$ -sum equal to zero. Hence there is no blue triangle.

Color green the remaining edges, i.e.  $k \ominus j \in \{1, 2\}$ . Difference=2 makes green square  $\{0, 2, 4, 6\}$ , and square  $\{1, 3, 5, 7\}$ . And difference=1 make a green octagon  $\{0, 1, 2, 3, 4, 5, 6, 7\}$ .

FTSOCContradiction, suppose  $G$  is a green  $K_4$  subgraph. It can't have three vertices in one square [diff=4], so  $G$  has two vertices in *each* square, say  $\{0, 2\}$  and  $\{v, w\}$ . But  $v$  must differ from each of 0,2 by 1 or 2; so  $v \overset{\text{must}}{=} 1$ . Ditto,  $w = 1$ .  $\bowtie$

The result now allows, courtesy Thm 2b, that

$$\mathcal{R}(4,4) \leq 2 \cdot \mathcal{R}(3,4) = 18.$$

3b: **Claim:**  $\mathcal{R}(4,4) = 18$ . ◆

**Proof.** To show  $\mathcal{R}(4,4) > 17$ , we exhibit a blue-green-coloring of  $K_{17}$  with no monochromatic  $K_4$ .

On vertex-set  $\mathbb{V} := \{0, 1, \dots, 8, -8, -7, \dots, -1\}$ , let  $\equiv, \oplus, \ominus$  and  $\langle \cdot \rangle$  each operate mod 17. Blue-connect  $j, k \in \mathbb{V}$  IFF  $k \ominus j$  is a mod-17 QR [Quadratic Residue].

3c:	$x$	$\langle x^2 \rangle$	$x$	$\langle x^2 \rangle$
	$\pm 1$	1	$\pm 5$	8
	$\pm 2$	4	$\pm 6$	2
	$\pm 3$	-8	$\pm 7$	-2
	$\pm 4$	-1	$\pm 8$	-4

Thus  $\text{QR} = \{\pm 1, \pm 2, \pm 4, \pm 8\}$ ; so our edge-lengths are  $\{1, 2, 4, 8\}$ . And  $\text{NQR} = \{\pm 3, \pm 5, \pm 6, \pm 7\}$ , giving rise to the green edges. Multiplying the vertices by a non-QR element, will exchange the blue and green edges [using that mult distributes-over addition]. So: *It suffices to show that there is no blue  $K_4$ .*

**Number Thy.** FTSOC, suppose  $\mathbf{a}, \mathbf{b}, \mathbf{c}, \mathbf{d} \in \mathbb{V}$  are the vertices of a blue  $K_4$ . WLOG  $\mathbf{a}=0$ ; replace vertex  $x$  by  $x \ominus \mathbf{a}$ . WLOG  $\mathbf{b}=1$ ; replace each  $x$  by  $\langle x/\mathbf{b} \rangle$ . [The mapping  $x \mapsto \langle x/\mathbf{b} \rangle$  is color-preserving, since  $\mathbf{b}$  is a QR.]

The vertices are now  $\{0, 1, \mathbf{c}, \mathbf{d}\}$ .

Vertex  $\mathbf{c}$  must differ from 0 and 1 by values in  $\{1, 2, 4, 8\}$ ; so  $\mathbf{c}$  is 2 or -1; WLOG  $\mathbf{c}=2$ . Vertex  $\mathbf{d}$ , relative to vertices 1 and 2, must be 3 or 0; but 0 is taken. Our supposedly-blue  $K_4$  is  $\{0, 1, 2, 3\}$ .

Alas, edge  $\overline{03}$  is green. ◆

## van der Waerden's thm

A *shift*  $s \in \mathbb{Z}$  and a *gap*  $g \in \mathbb{Z}_+$ , give rise to *finite* and *infinite arithmetic progressions*:

$$\begin{aligned} \mathbb{AP}_n(s, g) &:= \{s + jg \mid j \in [0..n)\}; \\ \mathbb{AP}(s, g) = \mathbb{AP}_\infty(s, g) &:= \{s + jg \mid j \in \mathbb{Z}\}. \end{aligned}$$

Use  $\mathbb{AP}_n$  to refer to *some*  $n$ -term A.P, when the shift and gap are irrelevant.

**4a: van der Waerden theorem.** *For each posint  $n$ , there is a finite  $M$  so that each 2-coloring (of the elements) of  $\mathbb{AP}_M$ , produces a monochromatic  $\mathbb{AP}_n$ .* ◇

The smallest such  $M$  is written  $\mathcal{W}(n)$ . For example,  $\mathcal{W}(3) = 9$ . If  $\mu$  many colors are allowed, then we write  $\mathcal{W}_\mu(n)$  for the minimum  $M$ .

The infinite analog of vdW thm fails, as shown next.

**4b: Obs.** *There exists a 2-coloring of  $\mathbb{Z}$  which admits no monochromatic sub- $\mathbb{AP}_\infty$ .* ◇

**Proof.** We produce a coloring for which the only monochromatic arithmetic-progs are finite.

Take a surjection  $n \mapsto (s_n, g_n)$  from  $\mathbb{Z}_+ \rightarrow \mathbb{Z} \times \mathbb{Z}_+$ , and let  $\mathcal{A}(n) := \mathbb{AP}_\infty(s_n, g_n)$ .

We color  $\mathbb{Z}$  in stages. At stage  $n$ : Paint one **blue** and one **green**, the two posints in  $\mathcal{A}(n)$  which are closest to zero, and were not yet colored. Then paint one **blue** and one **green**, the two negints in  $\mathcal{A}(n)$  which are closest to zero, and were not yet colored.

Send  $n \nearrow \infty$ . Finally, color arbitrarily the remaining uncolored integers. ◆

## Hypergraphs

Let  $\mathcal{P}_4(\mathbb{V})$  be the collection of cardinality-4 subsets of  $\mathbb{V}$ ; so it has  $\binom{|\mathbb{V}|}{4}$  many members. A “4-hypergraph” on  $\mathbb{V}$  is a pair  $H := (\mathbb{V}, \mathbb{E})$  where  $\mathbb{E} \subset \mathcal{P}_4(\mathbb{V})$  is the set of **hyper-edges** of  $H$ . This cardinality, 4, is called the **rank** of  $H$ .

A subset  $V' \subset \mathbb{V}$  determines a sub-hypergraph  $H' := (V', \mathbb{E}')$  of  $H$ , where

$$\mathbb{E}' := \{S \in \mathbb{E} \mid S \subset V'\}.$$

By the way, the “complete 4-hypergraph on  $\mathbb{V}$ ” is  $(\mathbb{V}, \mathcal{P}_4(\mathbb{V}))$ . Setting  $M := |\mathbb{V}|$ , we’ll use  $K_M^{(4)}$  to refer to this hypergraph. [So  $K_M^{(2)}$  is our usual  $K_M$ .]

5a: **Defn.** Fix a colorset  $\mathcal{C} := \{\text{blue, green}\}$ .

### A $\mathcal{C}$ -coloring of $H$

is a map  $\mathbf{f}: \mathbb{E} \rightarrow \mathcal{C}$ . Given a subset  $V' \subset \mathbb{V}$ ,

let  $\mathbf{f}|_{V'}$  mean the coloring  $\mathbf{f}|_{\mathbb{E}'}$  defined by (Y).

Consider a rank  $\rho \in [2.. \infty)$  and posint  $n$ . Suppose there is a posint  $M$  st:

For each  $\mathcal{C}$ -coloring of the complete hypergraph  $K_M^{(\rho)}$ , there is a cardinality- $n$  subset  $W \subset [1..M]$  so that coloring  $\mathbf{f}|_W$  is constant.

(IOWords, this edge-colored  $K_M^{(\rho)}$  admits a monochromatic  $K_n^{(\rho)}$ .) The smallest such  $M$  is the **hypergraph Ramsey number**  $R^{(\rho)}(n)$ .  $\square$

5b: **Hypergraph Ramsey Thm.** Fix a rank  $\rho \in \mathbb{Z}_+$ . Then

- i: Each coloring of  $K_\infty^{(\rho)}$  admits a monochromatic  $K_\infty^{(\rho)}$ -subgraph.
- ii: Each posint  $n$ : Ramsey number  $R^{(\rho)}(n)$  is finite.  $\diamond$

**Proof of (ii).** The analogous compactness argument of (1c) works here.  $\spadesuit$

**Proof of (i).** We induct on  $\rho$ . We’ll show the induction for  $\rho=5$ , assuming the  $\rho=4$  case.

Our vertex set is  $\mathbb{V} := \{1, 2, \dots\}$ , and we are given a coloring  $\mathbf{f}: \mathcal{P}_5(\mathbb{V}) \rightarrow \mathcal{C}$ . Suppose we could find an infinite subset  $\mathbf{W} \subset \mathbb{V}$  and a color map  $g: \mathcal{P}_4(\mathbf{W}) \rightarrow \mathcal{C}$  with this property:

†: For each  $S \in \mathcal{P}_4(\mathbf{W})$  and each  $y \in \mathbf{W}$  with  $y > \text{Max}(S)$ , the color  $\mathbf{f}(S \sqcup \{y\})$  equals  $g(S)$ .

The rank=4 case of (5b) asserts there is an  $\infty$ -subset  $\mathbf{X} \subset \mathbf{W}$  so that

Our  $g$ -coloring is constant on  $\mathcal{P}_4(\mathbf{X})$ ; say **blue**.

Given a  $T \in \mathcal{P}_5(\mathbf{X})$ , write it as  $T = \{w_1, \dots, w_4, w_5\}$  with  $w_1 < \dots < w_5$ . By (†), then,

$$\mathbf{f}(T) = g(\{w_1, \dots, w_4\}) = \text{blue}.$$

Hence  $\mathbf{f}$  is constant **blue** on  $\mathcal{P}_5(\mathbf{X})$ .

**Building  $\mathbf{W}$ .** We’ll inductively construct vertices  $w_1 < w_2 < \dots$  and infinite  $\mathbb{V}$ -subsets  $\mathbf{Y}_1 \supset \mathbf{Y}_2 \supset \dots$ . Our  $\mathbf{W}$  will be  $\{w_1, w_2, w_3, \dots\}$ .

Let  $w_1 := 1$ ,  $\mathbf{I}_1 := \{w_1\}$  and  $\mathbf{Y}_1 := [2.. \infty)$ .

At STAGE  $k$ : We have  $\mathbf{I}_k := \{w_1, \dots, w_k\}$ , and a partially-defined  $g()$ , defined on  $\mathcal{P}_4(\mathbf{I}_k)$ . We have an infinite vertex-set  $\mathbf{Y}_k$ , such that:

i: Our  $w_k < y$ , for each  $y \in \mathbf{Y}_k$ .

ii: For each  $S \in \mathcal{P}_4(\mathbf{I}_k)$ , and each  $y \in \mathbf{Y}_k$ , the color  $\mathbf{f}(S \sqcup \{y\})$  equals  $g(S)$ .

For STAGE  $[k+1]$ , define  $w_{k+1} := \text{Min}(\mathbf{Y}_k)$ , and temporary set

$$J_0 := \mathbf{Y}_k \setminus \{w_{k+1}\}.$$

Let  $S_1, S_2, \dots, S_L$  be some enumeration of those cardinality-4 subsets of  $\mathbf{I}_{k+1}$  that own  $w_{k+1}$ .

There is a color, say, **blue**, and an infinite set of  $y \in J_0$ , so that  $\mathbf{f}(S_1 \sqcup \{y\})$  is **blue**. Extend  $g()$  by defining  $g(S_1) := \text{blue}$ . Use  $J_1$  for this set of points  $y$ .

There is a color, say, **green**, and an infinite set of  $y \in J_1$ , so that  $\mathbf{f}(S_2 \sqcup \{y\})$  is **green**. Define  $g(S_2) := \text{green}$ . Use  $J_2$  for this set of points  $y$ .

Continue, until you have shrunk to  $J_L$ . Lastly, let  $\mathbf{Y}_{k+1} := J_L$ .  $\spadesuit$

5c: *Defn.* The **Erdős-Szekeres number**  $\text{ES}(n)$ , is the smallest posint  $M$  so that *each* collection of  $M$  points in the plane with no three colinear, has a subset of  $n$  points which form a *convex*  $n$ -gon. [Caveat: There are at least two different results called the Erdős-Szekeres thm.]  $\square$

5d: **ES-Theorem.** *For the Erdős-Szekeres number,*

$$\text{ES}(n) \leq R^{\langle 3 \rangle}(n, n) =: M. \quad \diamond$$

**Proof.** With vertex-set  $[1..M]$ , construct a **Cyan-Amber**-coloring of  $K_{[1..M]}^{\langle 3 \rangle}$ , as follows: For each triple  $\mathbf{u} < \mathbf{v} < \mathbf{w}$  in  $[1..M]$ , color the  $\{\mathbf{u}, \mathbf{v}, \mathbf{w}\}$ -edge **Cyan** if the  $\mathbf{u} \rightarrow \mathbf{v} \rightarrow \mathbf{w} \rightarrow \mathbf{u}$  traversal is ClockWise [**CW**]; otherwise, paint **Amber** the  $\{\mathbf{u}, \mathbf{v}, \mathbf{w}\}$ -edge, since its traversal is Anti-clockWise [**AW**].

By hypothesis, there exists an  $n$ -set  $S \subset [1..M]$ , so that all the  $K_S^{\langle 3 \rangle}$ -edges are, say, **CW**.

**Convex  $n$ -gon.** FTSOC, suppose the  $n$  points of  $S$  do *not* form a convex  $n$ -gon. Then some point  $P \in S$  is in the convex-hull of  $S$ . So there are distinct points  $\mathbf{u} < \mathbf{v} < \mathbf{w}$  in  $S$ , with an  $S$ -point  $P \in \text{Hull}(\{\mathbf{u}, \mathbf{v}, \mathbf{w}\})$ .

Recall  $\mathbf{u} \rightarrow \mathbf{v} \rightarrow \mathbf{w}$  is **CW**. We must have  $P > \mathbf{u}$ ; else  $P \rightarrow \mathbf{u} \rightarrow \mathbf{w}$  is **AW**.

And  $P^{\text{must}} > \mathbf{v}$ ; else  $\mathbf{u} \rightarrow P \rightarrow \mathbf{v}$  is **AW**. Continuing,  $P^{\text{must}} > \mathbf{w}$ ; else  $\mathbf{v} \rightarrow P \rightarrow \mathbf{w}$  is **AW**. But now,  $\mathbf{u} \rightarrow \mathbf{w} \rightarrow P$  is **AW**.  $\diamond$

## §A Hales-Jewett

### The Statement

Suppose we have a finite alphabet  $\mathbb{A} = \{a, b, \dots\}$  and we fix a length  $\mathbf{h}$ . A *degree- $\mathfrak{D}$  polynomial*  $f(x_1, \dots, x_{\mathfrak{D}})$  over  $\mathbb{A}^{\times \mathbf{h}}$  is a word

$$6: \quad f \in [\mathbb{A} \sqcup \{x_1, \dots, x_{\mathfrak{D}}\}]^{\times \mathbf{h}}$$

where each variable  $x_j$  occurs at least once in  $f$ . We evaluate  $f()$  at a  $\mathfrak{D}$ -tuple of  $\mathbb{A}$ -letters by plugging them in for the  $\mathfrak{D}$  variables. The range of this polynomial is a subset of  $\mathbb{A}^{\times \mathbf{h}}$  and has  $|\mathbb{A}|^{\mathfrak{D}}$  members. This range is called a  *$\mathfrak{D}$ -dimensional (affine) subspace*.

The RHS of (6) implies

$$7: \quad \begin{aligned} &\text{There are at most } [|\mathbb{A}| + \mathfrak{D}]^{\mathbf{h}} \text{ many} \\ &\mathfrak{D}\text{-dimensional subspaces of } \mathbb{A}^{\times \mathbf{h}}. \end{aligned}$$

By the way, a 1-dimensional subspace is also called an (affine) *line*.

The Hales-Jewett theorem states that given an *alphabet size*  $\alpha := |\mathbb{A}|$ , a number  $\mu$  of colors and a dimension  $\mathfrak{D}$ :

*There is a function  $\mathbf{h} = \mathbf{h}(\mathfrak{D}, \mu, \alpha)$  so that each  $\mu$ -coloring (coloring by  $\mu$  many colors) of the set of words  $\mathbb{A}^{\times \mathbf{h}}$  will have a monochromatic  $\mathfrak{D}$ -dimensional subspace.*

One cannot guarantee the stronger statement that there is a monochromatic subspace parallel to the coordinate axes ie. where each variable in the word of (6) occurs exactly once. This is already false in the  $\mathfrak{D}=1$  case: Let the *color* of a word in  $\{0, 1\}^{\times \mathbf{h}}$  be the mod-2 sum of its bits. Then a *line* consists of a pair of  $\mathbf{h}$ -words  $u0w$  and  $u1w$  differing in a single bit-position —which therefore have different colors.